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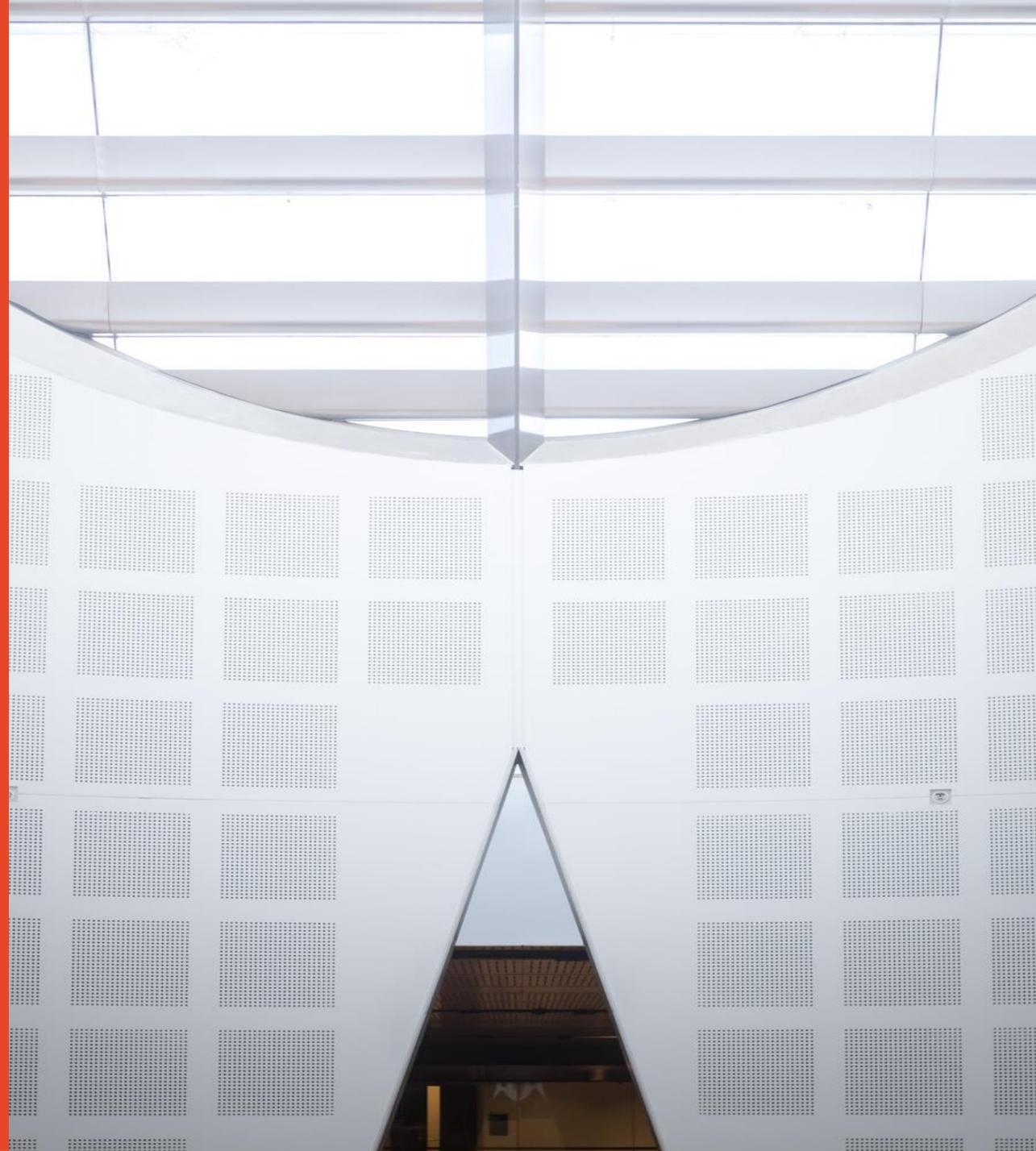
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COMPx270: Randomised and  
Advanced Algorithms  
Lecture 3: Balls in Bins

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SYDNEY



## A question 🎂

There are quite a few people in the classroom right now. What are the odds two of you (at least) have the same birthday?

Menti: 6529 7938



## A question

**Theorem.** (The  paradox) If you gather 23 people in a room, then with probability at least 50% a pair will sharing their birthday.

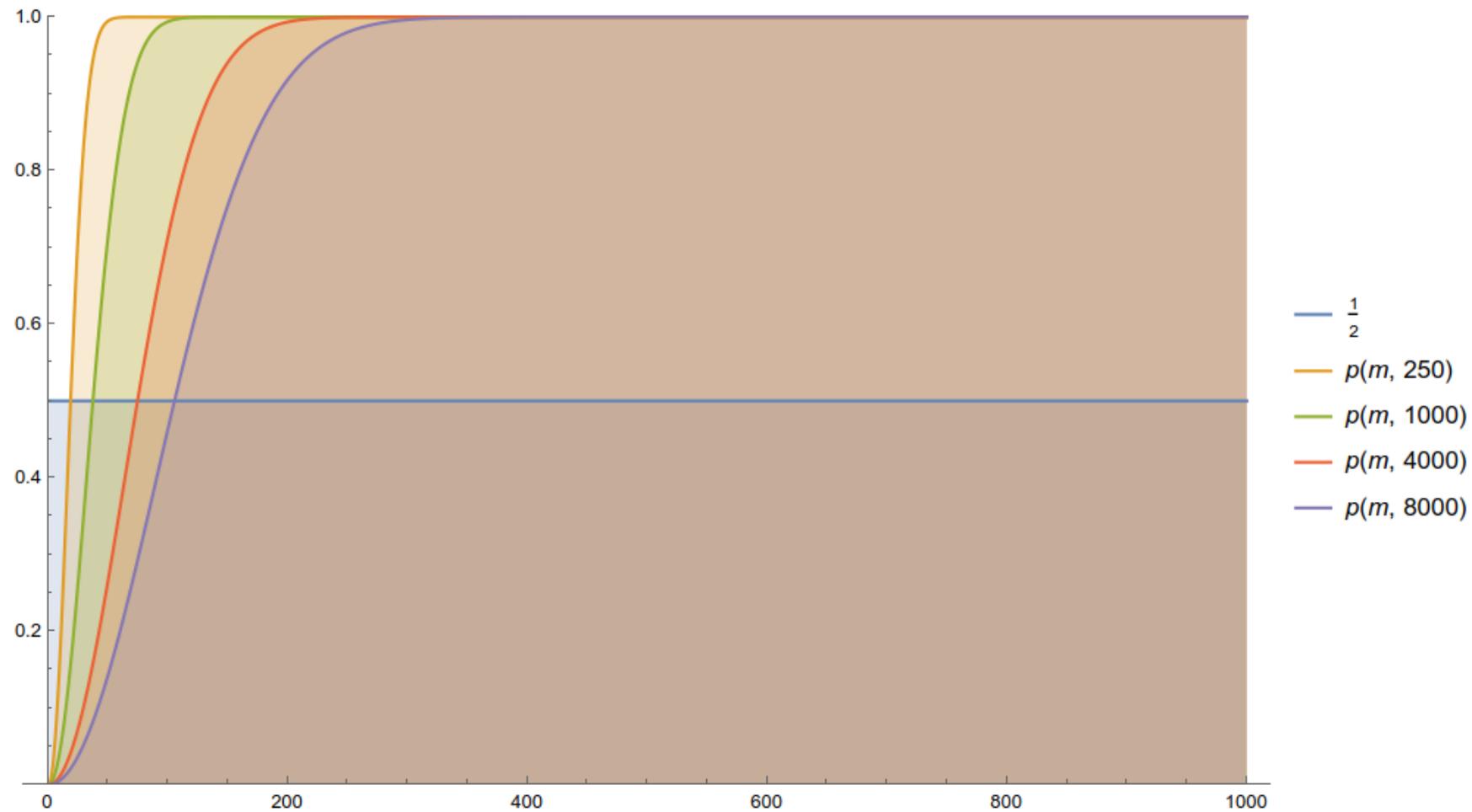
## An answer? 🎂

**Theorem.** If you gather  $m$  people and give each a number uniform between 1 and  $n$ , then the probability  $p(m, n)$  that at least two have the same number is...

$$p_{m,n} = 1 - \frac{n!}{n^m (n-m)!} = 1 - \frac{m!}{n^m} \binom{n}{m}$$

*Proof.*

```
DiscretePlot[{1/2, p[m, 250], p[m, 1000], p[m, 4000], p[m, 8000]}, {m, 1, 1000}, PlotRange -> {0, 1}, PlotLegends -> "Expressions"]
```



**Let's start simple:  $m=2$**   

**Now, for large values of "2"...**

C: number of collisions when throwing  $m$   into  $n$  . What is

$$c(m, n) = E[C]?$$

# Number of collisions

... and what is  $\text{Var}[C]$ ?



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for  $m = \Omega(\sqrt{n})$ . **Is it tight?**

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Now, we can use Chebyshev:

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for  $m = \Omega(\sqrt{n})$ .

By Markov, we also have

$$\Pr[ X \neq 0 ] = \Pr[ X \geq 1 ] \leq E[X] \leq 1/2$$

for  $m = O(\sqrt{n})$ .

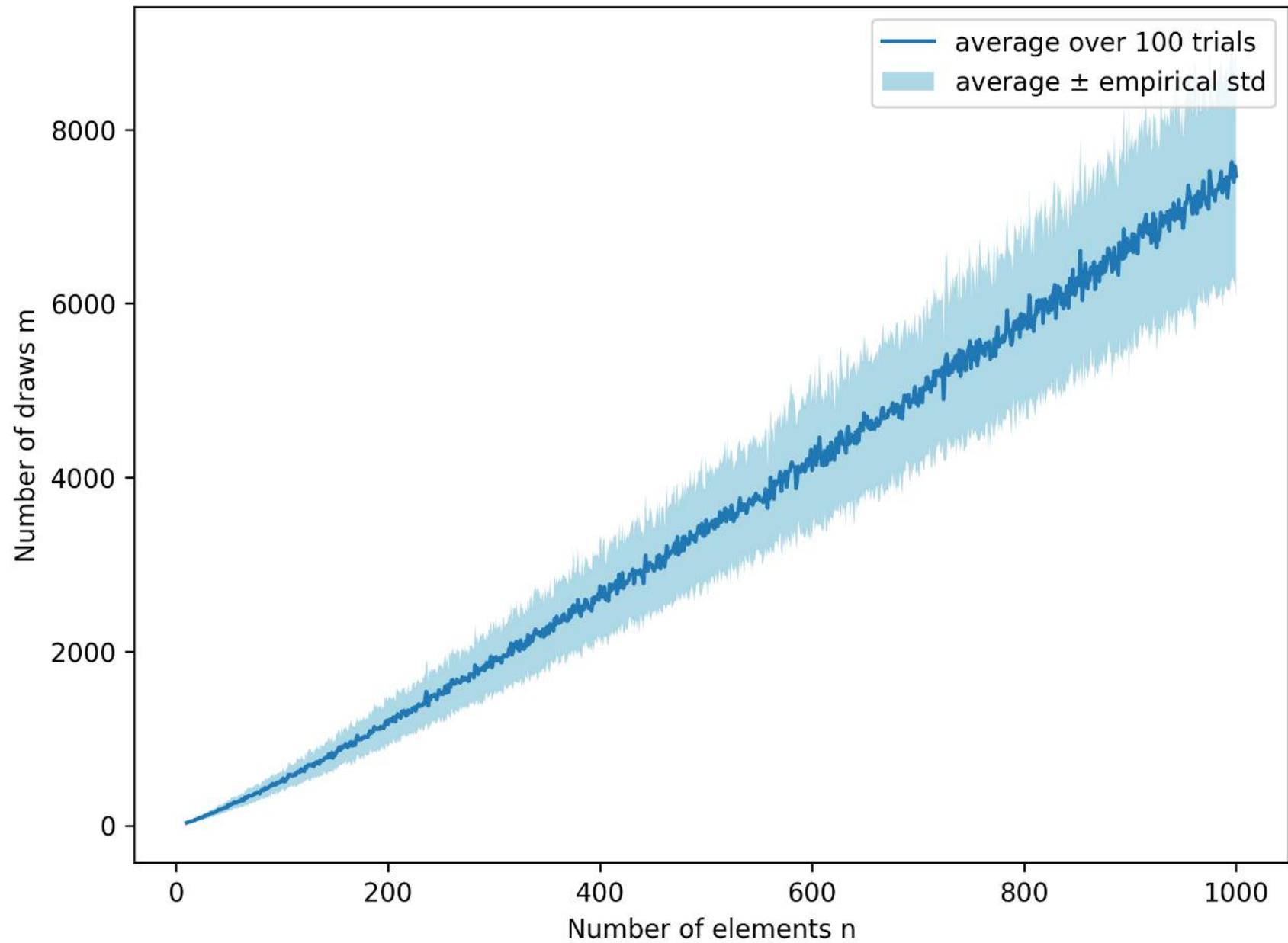
# Applications?

## **Bounding the variance: is it always that bad?**

Two tricks (and even 3).

# Coverage (Coupon collector)

"What is the expected number of balls  $M(n)$  we need to throw before each of the  $n$  bins contains at least one ball?"





# Load balancing

"What is the expected maximum number of balls  $L(n)$  any single of  $n$  bins contains when throwing  $n$  balls?"

